

Concurrency Control

These slides are a modified version of the slides of the book “Database System Concepts” (Chapter 16), 5th Ed., [McGraw-Hill](#), by Silberschatz, Korth and Sudarshan. Original slides are available at www.db-book.com

Lock-Based Protocols

Lock-Based Protocols

- A lock is a mechanism to control concurrent access to a data item
- Data items can be locked in two modes :
 1. *exclusive (X) mode*. Data item can be both read as well as written. X-lock is requested using **lock-X** instruction.
 2. *shared (S) mode*. Data item can only be read. S-lock is requested using **lock-S** instruction.
- Lock requests are made to concurrency-control manager. Transaction can proceed only after request is granted.

Lock-Based Protocols (Cont.)

- Lock-compatibility matrix

	S	X
S	true	false
X	false	false

- A transaction may be granted a lock on an item if the requested lock is compatible with locks already held on the item by other transactions
- Any number of transactions can hold shared locks on an item,
 - but if any transaction holds an exclusive on the item no other transaction may hold any lock on the item.
- If a lock cannot be granted, the requesting transaction is made to wait till all incompatible locks held by other transactions have been released. The lock is then granted.

Lock-Based Protocols (Cont.)

- Example of a transaction performing locking:

```
 $T_2$ : lock-S(A);  
      read (A);  
      unlock(A);  
      lock-S(B);  
      read (B);  
      unlock(B);  
      display(A+B)
```

- Locking as above is not sufficient to guarantee serializability — if A and B get updated in-between the read of A and B , the displayed sum would be wrong.
- A **locking protocol** is a set of rules followed by all transactions while requesting and releasing locks. Locking protocols restrict the set of possible schedules.

Pitfalls of Lock-Based Protocols

- Consider the partial schedule

T_3	T_4
lock-X(B)	
read(B)	
$B := B - 50$	
write(B)	
	lock-S(A)
	read(A)
	lock-S(B)
lock-X(A)	

- Neither T_3 nor T_4 can make progress — executing **lock-S(B)** causes T_4 to wait for T_3 to release its lock on B , while executing **lock-X(A)** causes T_3 to wait for T_4 to release its lock on A .
- Such a situation is called a **deadlock**.
 - To handle a deadlock one of T_3 or T_4 must be rolled back and its locks released.

Pitfalls of Lock-Based Protocols (Cont.)

- The potential for deadlock exists in most locking protocols. Deadlocks are a necessary evil.

- **Starvation** is also possible if concurrency control manager is badly designed. For example:
 - A transaction may be waiting for an X-lock on an item, while a sequence of other transactions request and are granted an S-lock on the same item.
 - The same transaction is repeatedly rolled back due to deadlocks.

- Concurrency control manager can be designed to prevent starvation.

The Two-Phase Locking Protocol

- This is a protocol which ensures conflict-serializable schedules.
- Phase 1: Growing Phase
 - transaction may obtain locks
 - transaction may not release locks
- Phase 2: Shrinking Phase
 - transaction may release locks
 - transaction may not obtain locks
- The protocol assures serializability. It can be proved that the transactions can be serialized in the order of their **lock points** (i.e. the point where a transaction acquired its final lock).

The Two-Phase Locking Protocol (Cont.)

- Two-phase locking *does not* ensure freedom from deadlocks
- Cascading roll-back is possible under two-phase locking. To avoid this, follow a modified protocol called **strict two-phase locking**. Here a transaction must hold all its exclusive locks till it commits/aborts.
- **Rigorous two-phase locking** is even stricter: here *all* locks are held till commit/abort. In this protocol transactions can be serialized in the order in which they commit.

The Two-Phase Locking Protocol (Cont.)

- There can be conflict serializable schedules that cannot be obtained if two-phase locking is used.

T1	T2	T3
read(x)		
write(x)		
	read(x)	
	write(x)	
		read(y)
write(y)		

Lock Conversions

- Two-phase locking with lock conversions:
 - First Phase:
 - can acquire a lock-S on item
 - can acquire a lock-X on item
 - can convert a lock-S to a lock-X (upgrade)
 - Second Phase:
 - can release a lock-S
 - can release a lock-X
 - can convert a lock-X to a lock-S (downgrade)
- This protocol assures serializability. But still relies on the programmer to insert the various locking instructions.

Automatic Acquisition of Locks

- A transaction T_i issues the standard read/write instruction, without explicit locking calls.
- The operation **read**(D) is processed as:

if T_i has a lock on D

then

read(D)

else begin

if necessary wait until no other

transaction has a **lock-X** on D

grant T_i a **lock-S** on D ;

read(D)

end

Automatic Acquisition of Locks (Cont.)

- **write**(D) is processed as:
 - if T_i has a **lock-X** on D
 - then**
 - write(D)
 - else begin**
 - if necessary wait until no other trans. has any lock on D ,
 - if T_i has a **lock-S** on D
 - then**
 - upgrade** lock on D to **lock-X**
 - else**
 - grant T_i a **lock-X** on D
 - write(D)
 - end;**
- All locks are released after commit or abort

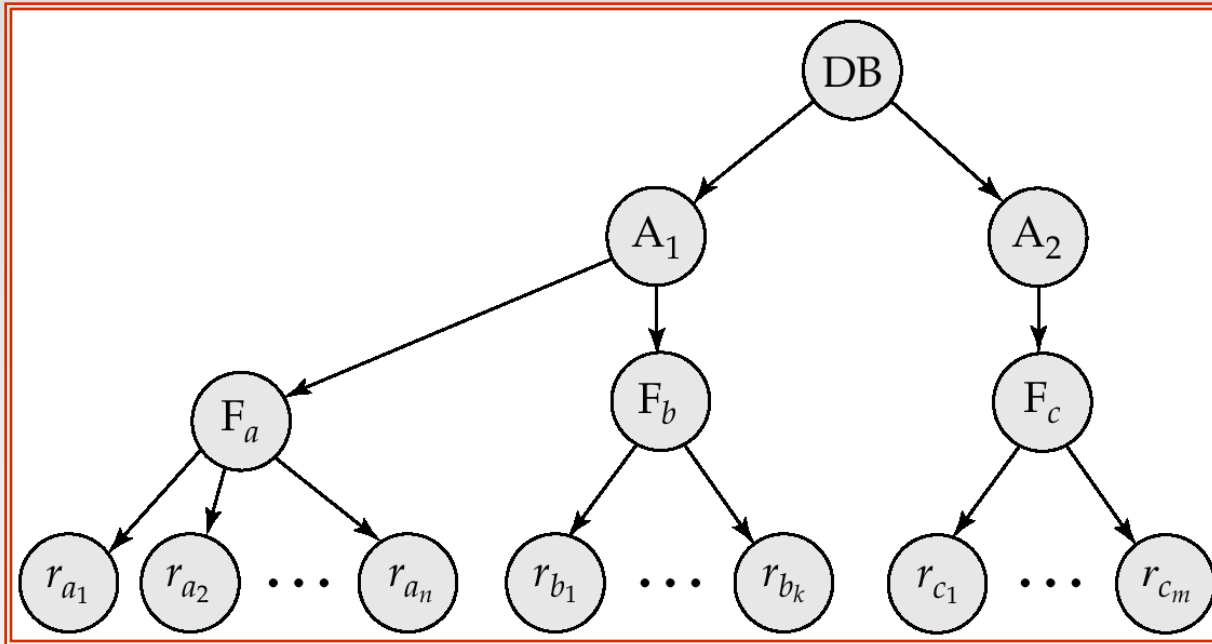
Implementation of Locking

- A **lock manager** can be implemented as a separate process to which transactions send lock and unlock requests
- The lock manager replies to a lock request by sending a lock grant messages (or a message asking the transaction to roll back, in case of a deadlock)
- The requesting transaction waits until its request is answered
- The lock manager maintains a data-structure called a **lock table** to record granted locks and pending requests

Multiple Granularity

- Allow data items to be of various sizes and define a hierarchy of data granularities, where the small granularities are nested within larger ones
- Can be represented graphically as a tree
- When a transaction locks a node in the tree *explicitly*, it *implicitly* locks all the node's descendents in the same mode.
- **Granularity of locking** (level in tree where locking is done):
 - **fine granularity** (lower in tree): high concurrency, high locking overhead
 - **coarse granularity** (higher in tree): low locking overhead, low concurrency

Example of Granularity Hierarchy



The levels, starting from the coarsest (top) level are

- *database*
- *area*
- *file*
- *record*

Intention Lock Modes

- In addition to S and X lock modes, there are three additional lock modes with multiple granularity:
 - ***intention-shared*** (IS): indicates explicit locking at a lower level of the tree but only with shared locks.
 - ***intention-exclusive*** (IX): indicates explicit locking at a lower level with exclusive or shared locks
 - ***shared and intention-exclusive*** (SIX): the subtree rooted by that node is locked explicitly in shared mode and explicit locking is being done at a lower level with exclusive-mode locks.

Compatibility Matrix with Intention Lock Modes

- The compatibility matrix for all lock modes is:

	IS	IX	S	S IX	X
IS	✓	✓	✓	✓	×
IX	✓	✓	×	×	×
S	✓	×	✓	×	×
S IX	✓	×	×	×	×
X	×	×	×	×	×

Multiple Granularity Locking Scheme

- Transaction T_i can lock a node Q , using the following rules:
 1. The lock compatibility matrix must be observed.
 2. The root of the tree must be locked first, and may be locked in any mode.
 3. A node Q can be locked by T_i in S or IS mode only if the parent of Q is currently locked by T_i in either IX or IS mode.
 4. A node Q can be locked by T_i in X, SIX, or IX mode only if the parent of Q is currently locked by T_i in either IX or SIX mode.
 5. T_i can lock a node only if it has not previously unlocked any node (that is, T_i is two-phase).
 6. T_i can unlock a node Q only if none of the children of Q are currently locked by T_i .
- Observe that locks are acquired in root-to-leaf order, whereas they are released in leaf-to-root order.

Deadlock Handling

- Consider the following two transactions:

T_1 : write (X)

write(Y)

T_2 : write(Y)

write(X)

- Schedule with deadlock

T_1	T_2
lock-X on X write (X) wait for lock-X on Y	lock-X on Y write (Y) wait for lock-X on X

Deadlock Handling

- System is deadlocked if there is a set of transactions such that every transaction in the set is waiting for another transaction in the set.
- **Deadlock prevention** protocols ensure that the system will *never* enter into a deadlock state. Some prevention strategies :
 - Require that each transaction locks all its data items before it begins execution (predeclaration).
 - Impose partial ordering of all data items and require that a transaction can lock data items only in the order specified by the partial order (graph-based protocol).

More Deadlock Prevention Strategies

- Following schemes use transaction timestamps for the sake of deadlock prevention alone.
- **wait-die** scheme — non-preemptive
 - older transaction may wait for younger one to release data item. Younger transactions never wait for older ones; they are rolled back instead.
 - a transaction may die several times before acquiring needed data item
- **wound-wait** scheme — preemptive
 - older transaction *wounds* (forces rollback) of younger transaction instead of waiting for it. Younger transactions may wait for older ones.
 - may be fewer rollbacks than *wait-die* scheme.

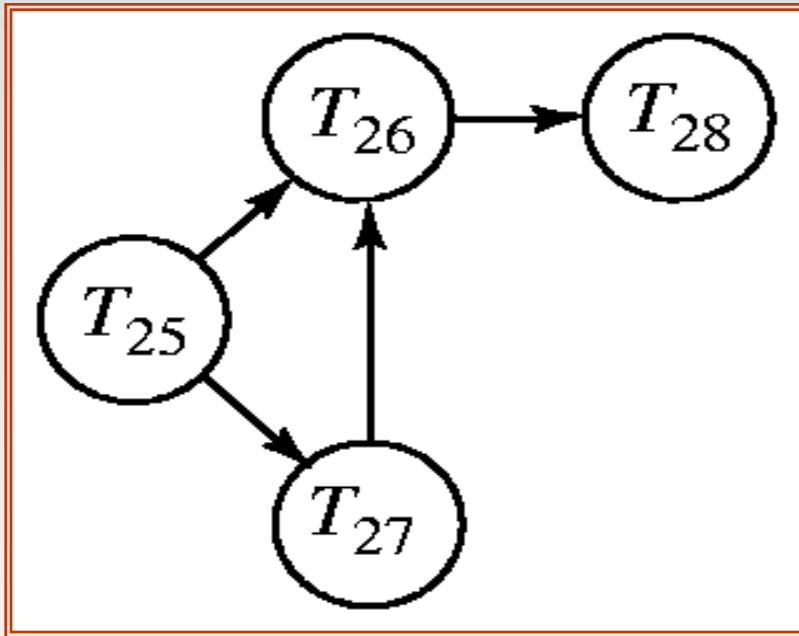
Deadlock prevention (Cont.)

- Both in *wait-die* and in *wound-wait* schemes, a rolled back transactions is restarted with its original timestamp. Older transactions thus have precedence over newer ones, and starvation is hence avoided.
- **Timeout-Based Schemes :**
 - a transaction waits for a lock only for a specified amount of time. After that, the wait times out and the transaction is rolled back.
 - thus deadlocks are not possible
 - simple to implement; but starvation is possible. Also difficult to determine good value of the timeout interval.

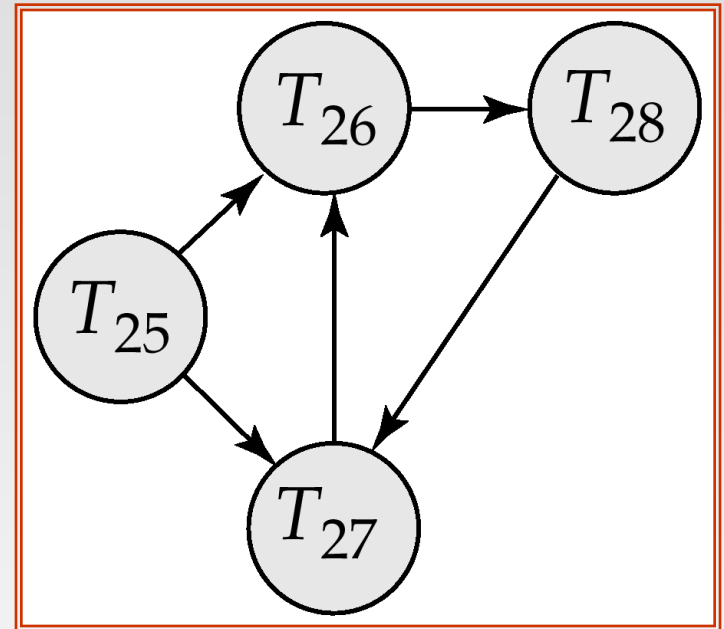
Deadlock Detection

- Deadlocks can be described as a *wait-for graph*, which consists of a pair $G = (V, E)$,
 - V is a set of vertices (all the transactions in the system)
 - E is a set of edges; each element is an ordered pair $T_i \rightarrow T_j$.
- If $T_i \rightarrow T_j$ is in E , then there is a directed edge from T_i to T_j , implying that T_i is waiting for T_j to release a data item.
- When T_i requests a data item currently being held by T_j , then the edge $T_i \rightarrow T_j$ is inserted in the wait-for graph. This edge is removed only when T_j is no longer holding a data item needed by T_i .
- The system is in a deadlock state if and only if the wait-for graph has a cycle. Must invoke a deadlock-detection algorithm periodically to look for cycles.

Deadlock Detection (Cont.)



Wait-for graph without a cycle



Wait-for graph with a cycle

Deadlock Recovery

- When deadlock is detected :
 - Some transaction will have to rolled back (made a victim) to break deadlock. Select that transaction as victim that will incur minimum cost.
 - Rollback -- determine how far to roll back transaction
 - ▶ **Total rollback**: Abort the transaction and then restart it.
 - ▶ More effective to roll back transaction only as far as necessary to break deadlock.
 - Starvation happens if same transaction is always chosen as victim. Include the number of rollbacks in the cost factor to avoid starvation

Timestamp-Based protocol

Timestamp-Based Protocols

- Each transaction is issued a timestamp when it enters the system. If an old transaction T_i has time-stamp $TS(T_i)$, a new transaction T_j is assigned time-stamp $TS(T_j)$ such that $TS(T_i) < TS(T_j)$.
- The protocol manages concurrent execution such that the time-stamps determine the serializability order.
- In order to assure such behavior, the protocol maintains for each data Q two timestamp values:
 - **W-timestamp**(Q) is the largest time-stamp of any transaction that executed **write**(Q) successfully.
 - **R-timestamp**(Q) is the largest time-stamp of any transaction that executed **read**(Q) successfully.

Timestamp-Based Protocols (Cont.)

- The timestamp ordering protocol ensures that any conflicting **read** and **write** operations are executed in timestamp order.
- Suppose a transaction T_i issues a **read**(Q)
 1. If $TS(T_i) \leq \mathbf{W}$ -timestamp(Q), then T_i needs to read a value of Q that was already overwritten.
 - Hence, the **read** operation is rejected, and T_i is rolled back.
 2. If $TS(T_i) \geq \mathbf{W}$ -timestamp(Q), then the **read** operation is executed, and R-timestamp(Q) is set to \mathbf{max} (R-timestamp(Q), $TS(T_i)$).

Timestamp-Based Protocols (Cont.)

- Suppose that transaction T_i issues **write**(Q).
 1. If $TS(T_i) < R\text{-timestamp}(Q)$, then the value of Q that T_i is producing was needed previously, and the system assumed that that value would never be produced.
 - Hence, the **write** operation is rejected, and T_i is rolled back.
 2. If $TS(T_i) < W\text{-timestamp}(Q)$, then T_i is attempting to write an obsolete value of Q .
 - Hence, this **write** operation is rejected, and T_i is rolled back.
 3. Otherwise, the **write** operation is executed, and $W\text{-timestamp}(Q)$ is set to $TS(T_i)$.

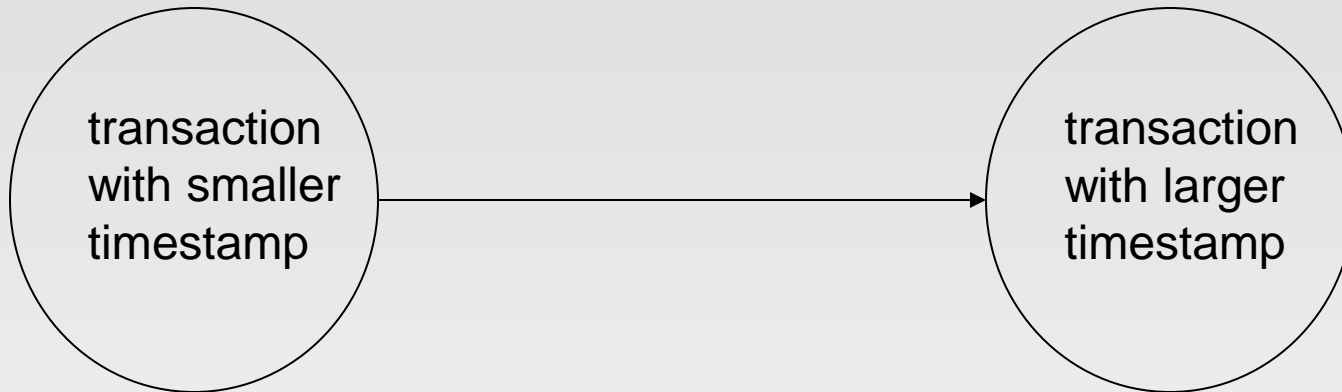
Example Use of the Protocol

A partial schedule for several data items for transactions with timestamps 1, 2, 3, 4, 5

T_1	T_2	T_3	T_4	T_5
read(Y)	read(Y)	write(Y) write(Z)		read(X)
read(X)	write(X) abort	write(Z) abort		read(Z)
				write(Y) write(Z)

Correctness of Timestamp-Ordering Protocol

- The timestamp-ordering protocol guarantees serializability since all the arcs in the precedence graph are of the form:



Thus, there will be no cycles in the precedence graph

- Timestamp protocol ensures freedom from deadlock as no transaction ever waits.
- But the schedule may not be cascade-free, and may not even be recoverable.

Timestamp-based Protocol

- There can be two-phase locking schedules that are not timestamp schedules

T1

T2

read(x)

write(x)

read(x)

write(x)

Recoverability and Cascade Freedom

- Problem with timestamp-ordering protocol:
 - Suppose T_i aborts, but T_j has read a data item written by T_i
 - Then T_j must abort; if T_j had been allowed to commit earlier, the schedule is not recoverable.
 - Further, any transaction that has read a data item written by T_j must abort
 - This can lead to cascading rollback --- that is, a chain of rollbacks
- Solution 1:
 - A transaction is structured such that its writes are all performed at the end of its processing
 - All writes of a transaction form an atomic action; no transaction may execute while a transaction is being written
 - A transaction that aborts is restarted with a new timestamp
- Solution 2: Limited form of locking: wait for data to be committed before reading it
- Solution 3: Use commit dependencies to ensure recoverability