#### DTMC: Discrete-time homogeneous Markov chain

- Markov process X (memoryless property) with finite discrete-state space S
- steady-state transition probabilities
- transitions at fixed intervals (steps)  $t \in \{0, 1, 2, ...\}$

#### Definitions

n number of states

 $\pi = (\pi_1, ..., \pi_n)$  state space distribution

Transition probability matrix

$$P = \begin{bmatrix} p_{11} \cdots p_{1n} \\ \vdots & \ddots \\ \vdots & \ddots \\ p_{n1} \cdots p_{nn} \end{bmatrix},$$

$$p_{ij} = \mathcal{P}\{X_1 = j | X_0 = i\}$$

probability of moving from state i to state j in one step

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Memoryless property

$$\mathcal{P}\{X_{t+1} = j | X_0 = k_0, ..., X_{t-1} = k_{t-1}, X_t = i\} = \mathcal{P}\{X_{t+1} = j | X_t = i\}$$

#### Characterization of time evolution of the process

 $p_{ij}^{(t)} = \mathcal{P}\{X_t = j | X_0 = i\}, t \in \{0, 1, 2, ...\}$ 

probability of moving from state i to state j in t steps

Transition probability from an initial state to a final state in (t1+t2) steps:

- transition probability from the initial state to a state k in t1 steps
- transition probability from state k to final state in t2 steps

$$p_{ij}^{(\texttt{t1+t2})} = \sum_{k \in \mathcal{S}} p_{ik}^{(\texttt{t1})} p_{kj}^{(\texttt{t2})} \; \forall i, j \in \mathcal{S} \; \forall \texttt{t1, t2} \geq 1$$

It follows that: t-steps Transition probability matrix

$$P^{(t)} = P^t$$
 *t-th power of P*

#### State space distribution

Initial state space distribution:  $\pi^{(0)} = (\pi_1^{(0)}, ..., \pi_n^{(0)})$ 

 $\pi_i^{(0)} = P\{X_0 = i\}$  initial probability vector

A single step forward:

$$\pi^{(1)} = \pi^{(0)} P$$

Probability that the DTMC is in state i at time-step t

$$\pi_i{}^{(t)} = \mathsf{P}\{X_t = i\}$$

#### Transient solution: $\pi^{(t)}$

State occupancy vector at time t in terms of the transition probability matrix:  $\pi^{(t)} = \pi^{(0)} \ P^t$ 

System evolution in a finite number of steps computed starting from the initial state distribution and the transition probability matrix

# Sojourn time

time spent by a DTMC in any of its states (independently of its initial distribution)

# ST<sub>i</sub>(k) probability that the DTMC stays in state i for k steps before moving to another state

Geometric distribution – random variable with memoryless property

Zi = number of steps that the DTMC stays in state i before moving  $S = \{0, 1\}$  0 stay in state i 1 move into a state different from i

$$P{Zi = k} = p_{ii}^{(k-1)} (1 - p_{ii})$$

 $p_{ii}$  probability of staying in state i (1- $p_{ii}$ ) probability of moving to another state

Z = number of trials before the first success (included the success)

$$ST_i(k) = P{Zi = k}$$

### Limiting behaviour

A DTMC can be specified in terms of the state occupancy probability vector  $\pi$  and a transition probability matrix P

$$\pi^{(t)} = \pi^{(0)} \mathsf{P}^t$$

The limiting behaviour of a DTMC (steady-state behaviour):

$$\lim_{\mathbf{t}\to\infty}\pi_j^{(\mathbf{t}\,)}$$

The limiting behaviour of a DTMC depends on the characteristics of its states. Sometimes the solution is simple.

## Irreducible DTMC

A state j is said to be **accessible** from state i if there exists t >=0 such that  $P_{ii}^{(t)} > 0$ , we write i->j

In terms of the graph, j is accessible from state i if there is a path from node i to node j.

A DTMC is **irreducible** if each state is accessible from every other state (for each i, j: it holds i ->j), otherwise it is **reducible** 

# Classification of states

A state i is **periodic** with period d > 1 if it is possible to move to state i only after n steps such that n = d, 2d, 3d, ...:

 $P_{ii}^{(t)} > 0$  implies t is an integer multiple of d

If d=1, the state i is said to be **aperiodic**; it is possible to move to the same state in one step



state 1 is periodic with period d=2 state 2 is periodic with period d=2

# Classification of states

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A state i is recurrent if
for each j: if (i->j) then (j->i)
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A state i is **transient** if exists (j!=i) such that (i->j) and not (j->i)

A state i is **absorbent** if p<sub>ii</sub>=1 (i is a recurrent state)

Each Markov chain has at least one recurrent state

### Steady-state behaviour

For aperiodic irreducible Markov chain for each j

$$\lim_{\mathbf{t}\to\infty}\pi_j^{(\mathbf{t}\,)}$$

exists and are independent from  $\pi^{(0)}$ 

If **all states are recurrent**, the **steady-state behaviour** of the Markov chain is given by the *fixpoint* of the equation:

$$\pi^{(t)} = \pi^{(t-1)} \mathsf{P}$$

with

$$\Sigma_j \pi_j = 1$$

 $\pi_i$  is inversely proportional to the period of recurrence of state j

#### Time-average state space distribution

For periodic Markov chains  $\lim_{\mathbf{t} \to \infty} \pi_j^{(\mathbf{t}\,)}$ 

doesn't exist (caused by the probability of the periodic state)

We compute the time-average state space distribution, called  $\pi^*$ 

$$\pi^* = \lim_{t \to \infty} \frac{\sum_{i=1}^{t} \pi^{(i)}}{\frac{1}{t}}$$

$$P = 1 \begin{bmatrix} 1 & 2 \\ 0 & 1 \\ 2 & 1 & 0 \end{bmatrix}$$



state i is periodic with period d=2

$$\begin{aligned} \pi^{(0)} &= (1,0) \\ \pi^{(1)} &= \pi^{(0)} P \\ \pi^{(2)} &= \pi^{(1)} P \\ \end{aligned} \qquad \begin{aligned} \pi^{(1)} &= (0,1) \\ \pi^{(2)} &= (1,0) \end{aligned}$$

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## Continuous-time homogeneous Markov chain (CTMC)

- Markov process X (memoryless property) with discrete-state space S
- steady-state transition rates
- events may occur at any point in time

#### Definitions

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Let T be an interval of real numbers (e.g., T=[0,1]). Memoryless property:

$$\mathcal{P}\{X_{t+\tau} = j | X_t = i, X_{t-t_1} = k_1, ..., X_{t-t_n} = k_n\} = \mathcal{P}\{X_{t+\tau} = j | X_t = i\}$$

= 0

for all  $\tau$ >0 and 0<t1<t2<...<tn.

Steady-state transition probabilities

$$\mathcal{P}\{X_{t+\tau} = j | X_t = i\} = \mathcal{P}\{X_\tau = j | X_0 = i\}$$

State-transition-rate matrix, the Q matrix

$$q_{ij} = \begin{cases} \text{rate of going from} & i \neq j, \\ \text{state } i \text{ to state } j & \sum_{j \in \mathcal{S}} q_{ij} \\ -\sum_{k \neq i} q_{ik} & i = j. \end{cases}$$

#### Single system with repair

- transition rates:  $\lambda$  failure rate,  $\mu$  repair rate
- identification of states
- initial state-space p(0) = [1, 0]

$$\begin{aligned} \pi^{(t)} &= \; [\pi_0^{(t)} \;, \; \pi_1^{(t)} \;] \\ p(t) &= [p_0(t), \; p_1(t)] \text{ in the book} \end{aligned}$$



$$\frac{dp_0(t)}{dt} = -\lambda p_0(t) + \mu p_1(t)$$
$$\frac{dp_1(t)}{dt} = \lambda p_0(t) - \mu p_1(t)$$

$$Q = \begin{bmatrix} -\lambda & \lambda \\ \mu & -\mu \end{bmatrix}$$

Solution of the differential equations:

$$p_0(t) = \frac{\mu}{\lambda + \mu} + \frac{\lambda}{\lambda + \mu} e^{-(\lambda + \mu)t}$$
$$p_1(t) = \frac{\lambda}{\lambda + \mu} - \frac{\lambda}{\lambda + \mu} e^{-(\lambda + \mu)t}$$

Availability  $A(t) = p_0(t)$ 

#### Transient solution

A CTMC can be specified in terms of the occupancy probability vector  $\pi$  and a transition probability matrix P

$$\pi^{(t)} = \pi^{(0)} \mathsf{P}^{t}$$

where

$$P^{(t)} = e^{Qt} \quad \text{for } t \ge 0$$

$$P^{(t)} = I + \frac{tQ}{1!} + \frac{t^2Q^2}{2!} + \frac{t^3Q^3}{3!} + \cdots$$

This allows to compute the probability of reaching state j from state i at time t :  $p_{ii}^{(t)}$ 

We have:

$$\pi^{(t)} = \pi^{(0)} e^{Qt}$$
$$\frac{d}{dt} \pi^{(t)} = \pi^{(t)} Q$$

Different numerical solution methods

# Sojourn time

For the memoryless property, the sojourn time spent by a CTMC in any of its states is independent of how long the CTMC has previously been in state i.

There is only one random variable that has this property: the exponential random variable:

ST<sub>i</sub> sojourn time in state i:

$$ST_i = e^{(a_i)}$$

-> the time spent in each state takes non-negative real values and has an exponential distribution

#### Steady-state behaviour

$$\lim_{t \to \infty} \pi_j^{(t)}, \, \forall j \in \mathcal{S}$$

We can prove that we have to solve the equation:

$$\pi^* \mathsf{Q} = \mathsf{0}, \quad ext{ where } \pi^* = \lim_{t o \infty} \pi_j^{(t)}$$

For **irreducible CTMC** (irreducible if every state is accessible from every other state. For each i, j, it holds i ->j.) the solution can be calculated under the constraint:

$$\sum_{i=1}^n \pi_i^* = 1$$

The steady-state distribution is independent of the initial-state distribution

If the CTMC is not **irreducible** then more complex solution methods are required

#### Direct methods:

Good packages exists Very poor performance if Q is very large

#### Iterative methods:

An iterative method converges if :

$$\lim_{k \to \infty} \left\| \pi^{(k)} - \pi^* \right\| = 0$$
  
Stopping condition: 
$$\left\| \pi^{(k+1)} - \pi^{(k)} \right\| < \varepsilon$$

Other methods

## Dual processor system with repair

#### A, B processors



#### **Reliability modeling**

- making state 2 a trapping state

p(0) = [1, 0, 0]



**Reliability**  $R(t) = 1 - p_2(t)$   $R(t) = p_0(t) + p_1(t)$ Laplace transform

$$R(t) = \frac{4\lambda^{2} \exp(-(1/2)(3\lambda + \mu - \sqrt{\lambda^{2} + 6\lambda\mu + \mu^{2}})t)}{(3\lambda + \mu)\sqrt{\lambda^{2} + 6\lambda\mu + \mu^{2}} - \lambda^{2} - 6\lambda\mu - \mu^{2}} - \frac{4\lambda^{2} \exp(-(1/2)(3\lambda + \mu + \sqrt{\lambda^{2} + 6\lambda\mu + \mu^{2}})t)}{(3\lambda + \mu)\sqrt{\lambda^{2} + 6\lambda\mu + \mu^{2}} + \lambda^{2} + 6\lambda\mu + \mu^{2}}$$

# TMR system with repair

Rates:  $\lambda$  and  $\mu$ 

Identification of states:

- 3 processors working, 0 failed
- 2 processors working, 1 failed
- 1 processor working, 2 failed



Transition rate matrix:

$$Q = \begin{bmatrix} -3\lambda & 3\lambda & 0\\ \mu & -2\lambda - \mu & 2\lambda\\ 0 & 0 & 0 \end{bmatrix} \qquad P(0) = \begin{bmatrix} 1, 0, 0 \end{bmatrix}$$

**Reliability** R(t) = 1 - p2(t) Laplace transform

$$R(t) = \frac{5\lambda + \mu + \sqrt{\lambda^2 + 10\lambda\mu + \mu^2}}{2\sqrt{\lambda^2 + 10\lambda\mu + \mu^2}} \exp(-(1/2)(5\lambda + \mu - \sqrt{\lambda^2 + 10\lambda\mu + \mu^2})t)$$
$$- \frac{5\lambda + \mu - \sqrt{\lambda^2 + 10\lambda\mu + \mu^2}}{2\sqrt{\lambda^2 + 10\lambda\mu + \mu^2}} \exp(-(1/2)(5\lambda + \mu + \sqrt{\lambda^2 + 10\lambda\mu + \mu^2})t)$$

# Comparison with nonredundant system and TMR without repair



$$\mathsf{MTTF} = \int_{t=0}^{\infty} \mathsf{R}(\mathsf{t}) \, \mathsf{d}\mathsf{t}$$

period the system is in a state that correspond to correct behavior

TMR with repair:

MTTF = 
$$\int_{t=0}^{\infty} p_0(t) + p_1(t) dt$$

failure rate  $\lambda = 0.001$  repair rate  $\mu = 0.1$ 

TMR with repair MTTF =  $\frac{5}{6\lambda} + \frac{\mu}{6\lambda^2} = 17,5000$  hours

MTTF is equal to the MTTF of a TMR system without repair plus an additional term due to the repair activity.

Nonredundant MTTF = 
$$\frac{1}{\lambda}$$
 = 1000 hours  
TMR without repair MTTF =  $\frac{5}{6\lambda}$  = 833 hours

on-line repair allows the system MTTF to increase by a factor of 17

## An example of modeling (CTMC)

Multiprocessor system with 2 processors and 3 shared memories system. System is operational if at least one processor and one memory are operational.



 $\lambda_m$  failure rate for memory  $\lambda_p$  failure rate for processor

X random process that represents the number of operational memories and the number of operational processors at time t

Given a state (i, j): i is the number of operational memories; j is the number of operational processors

 $S = \{(3,2), (3,1), (3,0), (2,2), (2,1), (2,0), (1,2), (1,1), (1,0), (0,2), (0,1)\}$ 

## Reliability modeling



 $\lambda_m$  failure rate for memory  $\lambda_p$  failure rate for processor

 $(3, 2) \rightarrow (2, 2)$  failure of one memory

(3,0), (2,0), (1,0), (0,2), (0,1) are absorbent states

## Availability modeling

- Assume that faulty components are replaced and we evaluate the probability that the system is operational at time t
- > Constant repair rate  $\mu$  (number of expected repairs in a unit of time)
- Strategy of repair: only one processor or one memory at a time can be substituted
- The behaviour of components (with respect of being operational or failed) is not independent: it depends on whether or not other components are in a failure state.

#### Strategy of repair:

only one component can be substituted at a time



 $\lambda$ m failure rate for memory  $\lambda$ p failure rate for processor  $\mu$ m repair rate for memory  $\mu$ p repair rate for processor An alternative strategy of repair:

only one component can be substituted at a time and processors have higher priority

exclude the lines µm representing memory repair in the case where there has been a process failure



## System model analysis

What is the availability of the system at time t?

What is the steady-state availability?

What is the expected time to failure?

The Markov model fits with the standard assumption of failure rates a constant, leading to exponentially distributed inter-arrival times of failures. Similarly, we assume costant **repair rate**.

What about safety?

## Safety

Safety - avoidance of catastrophic consequences -As a function of time, S(t), is the probability that the system either behaves correctly or will discontinue its functions in a manner that causes no harm (operational or Fail-safe)

**Coverage** – The coverage is the measure **c** of the system ability to reach a fail-safe state after a fault.

Modeling coverage and safety in a Markov chain means that every unfailed state has two transitions to two different states, one of which is fail-safe, the other is fail-unsafe.



# TMR

the system can be in a safe state although the failures of two components, if the output of the three components disagree

c = probability of coincident failures of two components



- 0 three correct components
- 1 one faulty component
- 2 two faulty components (no coincident failures)
- 3 two faulty component coincident failures
- 4 three faulty components (no coincident failures)

## Observations

Quantitative dependability evaluation:

- guiding design decisions
- assessing systems as built
- mandatory for safety critical systems

Model construction techniques

-> scalability challenge

#### composition approaches

build complex models in a modular way through a composition of its submodels

#### decomposition/aggregation approaches

(hierarchical decomposition approach)

The overall model is decoupled in simpler and more tractable submodels, and the measures obtained from the solution of the submodels are then aggregated to compute those concerning the overall model.