

Exercise 1 (9 points)

Let's consider the following relational schema for an insurance company, customers and policies:

CUSTOMER(Cust_Id, Name, Age, City)

BRANCH(Branch_Id, Branch_name, nEmployee, Branch_city)

POLICY(Policy_number, Cust_Id, Branch_Id, expiry_date)

Primary keys are underlined in the relations. Moreover, Cust_Id in POLICY is foreign key of CUSTOMER and Branch_Id in POLICY is foreign key of BRANCH.

A customer can have more than one policy in the same or in different branches.

Assume that:

$n_{\text{CUSTOMER}} = 50.000$

$n_{\text{BRANCH}} = 10$

$n_{\text{POLICY}} = 100.000$

$V(\text{Cust_Id}, \text{POLICY}) = 50.000$

$V(\text{Branch_Id}, \text{POLICY}) = 10$

$V(\text{City}, \text{CUSTOMER}) = 15$

$V(\text{Branch_city}, \text{BRANCH}) = 5$

$V(\text{expiry_date}, \text{POLICY}) = 20$

Given the query:

"Branch name of branches located in Florence holding policies with expiry date = 2018 stipulated by customers leaving in Pisa."

- 1) express the query as a relational-algebra expression;
- 2) show the basic steps of the query optimization process in terms of relational-algebra expression transformations
- 3) give an efficient strategy for computing the query.

Exercise 2 (6 points)

Consider the following schedule of concurrent transactions:

S: r1(x) r1(v) r3(z) w2(z) r3(x) w3(y) w1(y) w2(v)

- 1) Show if it is conflict serializable (CSR) or view serializable (VSR). Explain why. If serializable, show equivalent serial schedules.
- 2) Apply the rigorous two-phase locking protocol to the schedule. Is the schedule accepted?
- 3) Apply the timestamp-ordering protocol to the schedule, assuming that aborted transactions are immediately restarted. Is the schedule accepted?

Exercise 3 (6 points)

Consider an empty B+-tree with $m=4$.

1) Show the B+-tree after the insertion of the following values of the search key:

36 12 18 25 3 70 19 5 40 30

2) Show the form of the B+-tree after the following operation:

Delete 70.

Exercise 4 (9 points)

Let $r=(A,B,C)$, with A a key.

Assume

$nr = 1.000.000$	number of records in the relation
$Lr = 50$ byte	size of a record (fixed length records)
$LA = 6$ byte	size of attribute A
$Lp = 4$ byte	size of a pointer
$Lb = 1000$ byte	size of a block

Heap file organization.

1. Show the number of blocks of a **B+-tree index** on search-key A, assuming that each block contains 70% of the records that can fit in a block.

2. Outline the steps in answering the following queries, showing **the best strategy** and **the cost** in terms of number of block transfers from disk:

1) select * from r where $A=xxx$;

2) select * from r where $50.000 \leq A < 250.000$;
assuming A uniformly distributed on the interval $[1; 10.000.000]$

3) select * from R where $B=xxx$;