## **Exercise 1**

Consider the following function of two variables:

$$
f(x,y) = \begin{cases} C \cdot (x+y) & x \in [0,1], \ y \in [0,1], 0 < x+y < 1 \\ 0 & otherwise \end{cases}
$$

- 1) Compute C so that  $f(x, y)$  is a JPDF for RVs *X* and *Y*
- 2) Compute the PDFs for RVs  $X$  and  $Y$  and compute  $Cov(X, Y)$
- 3) Define  $Z = X + Y$ . Compute the PDF of *Z* and determine  $E[Z]$

## **Exercise 2**

A server's main job is to handle transactions. These arrive exponentially at a rate  $\lambda$ , they are queued into an infinite FIFO buffer, and they are served at an exponential rate  $\mu$ . However, transaction handling is sometimes interrupted by *garbage collection (GC) requests*. When one GC request arrives, the sever *stops accepting and serving transactions* and attends to the GC. When the GC request has been cleared, the server resumes accepting/serving transactions. GC requests arrive at a rate  $\gamma$  and are served at a rate  $\delta$ . However, a GC request may arrive only while the server is handling transactions, and there can only be one outstanding GC request (i.e., GC requests cannot queue up).

- 1) Model the system and draw the CTMC
- 2) Write the steady-state equations, find the stability condition and compute the SS probabilities. Explain the stability condition.
- 3) Express the fraction of time that the server spends handling GC requests.
- 4) Compute the mean number of transactions in the system.
- 5) Compute the transaction throughput.
- 6) Determine what happens of 3,4,5 in the limit cases  $\gamma \ll \delta$  and  $\gamma \gg \delta$ .

## **Exercise 1 – Solution**

The JPDF is non null in the triangle in the figure.



Therefore, normalization reads:

$$
\int_0^1 \left[ \int_0^{1-x} C(x+y) dy \right] dx = C \cdot \int_0^1 \left[ \left[ x \cdot y + \frac{y^2}{2} \right]_0^{1-x} \right] dx = C \cdot \int_0^1 \left( \frac{x}{2} - x^2 + \frac{1+x^2 - 2x}{2} \right) dx = \frac{C}{2}.
$$
  

$$
\left[ x - \frac{x^3}{3} \right]_0^1 = \frac{C}{3} = 1
$$

Which yields  $C = 3$ .

2) To obtain the PDF for a RV, we integrate in the other.

$$
f_X(x) = \int_0^{1-y} 3(x+y) dy = \frac{3}{2} (1-x^2)
$$
, with  $x \in [0,1]$ .

For obvious reasons of symmetry,  $f_Y(y) = \frac{3}{2}$  $\frac{3}{2}(1-y^2)$ , with  $y \in [0,1]$ .

It is 
$$
Cov(X, Y) = E[XY] - E[X]E[Y]
$$
. Therefore:  
\n
$$
E[XY] = \int_0^1 \left[ \int_0^{1-x} 3xy(x+y) dy \right] dx = 3 \int_0^1 \left[ \int_0^{1-x} (x^2 \cdot y + x \cdot y^2) dy \right] dx
$$
\n
$$
= 3 \int_0^1 \left[ x^2 \cdot \frac{1+x^2 - 2x}{2} + x \cdot \frac{1-3x+3x^2 - x^3}{3} \right] dx = \frac{1}{2} \left[ x^2 - x^3 + \frac{x^5}{5} \right]_0^1 = \frac{1}{10}
$$

 $E[X] = E[Y] = \int_0^1 x \cdot \frac{3}{2}$  $\frac{3}{2}(1-x^2)dx=\frac{3}{8}$ 8 1 0

Therefore,  $Cov(X, Y) = E[XY] - E[X]E[Y] = \frac{1}{10}$  $\frac{1}{10} - \frac{9}{64}$  $\frac{9}{64} = -\frac{13}{320}$ 320

3) By the linearity of the mean value, one can immediately write  $E[X] = E[Z] + E[Y] = \frac{6}{8}$  $\frac{6}{8}$ 

The PDF of RV *Z* is non null in [0,1], and  $F_Z(z) = P\{X + Y \leq z\}$  is the probability that X and Y are in the triangle bounded by the dashed line in the figure.



Therefore, computing  $F_Z(z)$  only entails substituting  $z$  for 1 in the computations already done for the normalization:

$$
F_Z(z) = \int_0^z \left[ \int_0^{z-x} 3(x+y) dy \right] dx = 3 \cdot \int_0^z \left[ \left[ x \cdot y + \frac{y^2}{2} \right]_0^{z-x} \right] dx
$$
  
=  $3 \cdot \int_0^z \left( x \cdot z - x^2 + \frac{z^2 + x^2 - 2z \cdot x}{2} \right) dx = \frac{3}{2} \cdot \left[ z^2 \cdot x - \frac{x^3}{3} \right]_0^z = \frac{3}{2} \cdot \frac{2z^3}{3} = z^3$ 

From the above, we straightforwardly obtain  $f_Z(z) = 3z^2$ .

## **Exercise 2 – solution**

The CTMC is the following:



Calling  $P_j$  the probabilities of the "upper" states and  $P_j'$  those of "lower" states, the SS equations are the following:

$$
P_0 \cdot \lambda = P_1 \cdot \mu
$$
  

$$
P_j \cdot (\lambda + \mu) = P_{j+1} \cdot \mu + P_{j-1} \cdot \lambda, \quad j > 0
$$
  

$$
P'_j \cdot \delta = P_j \cdot \gamma, \quad j > 0
$$

From these, one obtains that  $P_j = P_0 \cdot \left(\frac{\lambda}{\mu}\right)$  $\left(\frac{\lambda}{\mu}\right)^j$ , and the following normalization equation:

$$
P_0 \cdot \left[ 1 + \left( 1 + \frac{\gamma}{\delta} \right) \cdot \sum_{j=1}^{+\infty} \left( \frac{\lambda}{\mu} \right)^j \right] = 1
$$

From which the stability condition is  $\lambda < \mu$ , and it is independent of  $\lambda$  or  $\delta$ . This is expectable, since during GC neither arrivals nor services occur, hence the transaction queue does not build up. Calling  $\rho = \lambda/\mu$ , and  $\theta = \lambda/\delta$ , we get:

$$
P_0 = \tfrac{1}{1+(1+\theta)\cdot\frac{\rho}{1-\rho}} = \tfrac{1-\rho}{1+\theta\cdot\rho}, \quad P_j = \tfrac{1-\rho}{1+\theta\cdot\rho}\cdot\rho^j, \quad P_j' = \tfrac{1-\rho}{1+\theta\cdot\rho}\cdot\theta\cdot\rho^j
$$

The fraction of time that the server spends attending GC requests is

$$
P_{GC} = \sum_{j=1}^{+\infty} P'_j = \frac{1-\rho}{1+\theta \cdot \rho} \cdot \theta \cdot \sum_{j=1}^{+\infty} \rho^j = \frac{\theta \cdot \rho}{1+\theta \cdot \rho}
$$

The mean number of transactions in the system is

$$
E[N] = \sum_{j=1}^{+\infty} j \cdot (P_j + P_j') = \frac{1 - \rho}{1 + \theta \cdot \rho} \cdot (1 + \theta) \cdot \sum_{j=1}^{+\infty} j \cdot \rho^j = \frac{\rho}{1 - \rho} \cdot \frac{1 + \theta}{1 + \theta \cdot \rho}
$$

The throughput is:

$$
tpt = \sum_{j=1}^{+\infty} \mu_j \cdot P_j = \mu \cdot \frac{1-\rho}{1+\theta \cdot \rho} \cdot \sum_{j=1}^{+\infty} \rho^j = \mu \cdot \frac{1-\rho}{1+\theta \cdot \rho} \cdot \frac{\rho}{1-\rho} = \frac{\lambda}{1+\theta \cdot \rho}
$$

The limit cases  $\gamma \ll \delta$  and  $\gamma \gg \delta$  correspond to  $\theta \cong 0$  and  $\theta \to \infty$ .

- a)  $\theta \cong 0$ :  $P_{GC} \cong 0$ ,  $E[N] = \frac{\rho}{1-\rho}$  $\frac{\rho}{1-\rho}$ , tpt =  $\lambda$ . The system behaves like an M/M/1, which is obvious since GCs cause negligible overhead.
- b)  $\theta \to \infty$ :  $P_{GC} \to 1$ ,  $E[N] \to \frac{1}{1-\theta}$  $\frac{1}{1-\rho}$ , t $pt \to 0$ . The system throughput goes down to zero. The transaction queue, however, does not grow indefinitely.