Exercise 1

Given an exponential RV Y, whose rate is α , consider RV $X = x_m \cdot e^Y$, $x_m > 0$.

- 1) Compute the interval of values for *X*, its CDF and its PDF
- 2) Compute the mean value of X and its variance
- 3) Is X heavy-tailed? Justify your answer
- 4) Compute conditional probability $P\{X > a + b | X > a\}$, with *a,b* positive values. Is the distribution memoryless? What happens to the above probability if $a \to +\infty$?

Exercise 2

Consider a computer system to which jobs are sent through an encrypted channel, with exponential interarrival times at a rate λ . When a *new* job arrives, the system first *decrypts* it. Decryption takes an exponential time, with a mean $\frac{1}{\alpha}$. Decrypted jobs are served in FIFO order with an exponential rate equal to μ . However, the system *cannot* perform any other task when decrypting a new job: either it is accepting/serving jobs, or it is decrypting one. More specifically, any arrival that occurs when the system is decrypting a job is lost, and any ongoing service is suspended and then resumed afterwards.

- 1) Model the system as a queueing system and draw the transition-rate diagram
- 2) Find the stability condition and compute the steady-state probabilities. State explicitly their dependence on α , and justify your answer.
- 3) Compute the mean arrival rate at the system and the tagged-arrival probabilities.
- 4) Compute the mean number of jobs in the system and the mean response time.

Exercise 1 – Solution

1) The interval of values for *X* is x_m , $+\infty$).

$$F_X(x) = P\{X \le x\} = P\{x_m \cdot e^Y \le x\} = P\left\{Y \le \log\left(\frac{x}{x_m}\right)\right\} = 1 - e^{-\alpha \cdot \log\left(\frac{x}{x_m}\right)} = 1 - \left(\frac{x_m}{x}\right)^{\alpha}$$
$$f_X(x) = \frac{\partial}{\partial x} F_X(x) = \frac{\alpha \cdot x_m^{\alpha}}{x^{\alpha+1}}$$

2) The first two moments are:

$$E[X] = \int_{x_m}^{+\infty} x \cdot \frac{\alpha \cdot x_m^{\alpha}}{x^{\alpha+1}} dx = \frac{\alpha \cdot x_m^{\alpha}}{1-\alpha} \cdot \left[x^{(1-\alpha)}\right]_{x_m}^{+\infty} = \begin{cases} \frac{\alpha}{\alpha-1} \cdot x_m & \alpha > 1\\ +\infty & otherwise \end{cases}$$
$$E[X^2] = \int_{x_m}^{+\infty} x^2 \cdot \frac{\alpha \cdot x_m^{\alpha}}{x^{\alpha+1}} dx == \alpha \cdot x_m^{\alpha} \int_{x_m}^{+\infty} x^{-\alpha+1} dx = \frac{\alpha \cdot x_m^{\alpha}}{2-\alpha} \cdot \left[x^{(2-\alpha)}\right]_{x_m}^{+\infty}$$
$$= \begin{cases} \frac{\alpha}{\alpha-2} \cdot x_m^2 & \alpha > 2\\ +\infty & otherwise \end{cases}$$

From the above, we obtain:

$$\sigma^{2} = E[X^{2}] - E[X]^{2} = \begin{cases} \frac{\alpha}{\alpha - 2} \cdot x_{m}^{2} - \frac{\alpha^{2}}{(\alpha - 1)^{2}} \cdot x_{m}^{2} & \alpha > 2 \\ +\infty & otherwise \end{cases}$$
$$= \begin{cases} \frac{\alpha}{(\alpha - 2)} \cdot \left(\frac{x_{m}}{\alpha - 1}\right)^{2} & \alpha > 2 \\ +\infty & otherwise \end{cases}$$

3) The definition of heavy-tailed distribution is: $\forall \lambda > 0$, $\lim_{x \to \infty} e^{\lambda x} \cdot (1 - F(x)) = \infty$ In this case, we have $\forall \lambda > 0$, $\lim_{x \to \infty} e^{\lambda x} \cdot (\frac{x_m}{x})^{\alpha} = +\infty$, hence the above distribution is heavy-tailed.

4) The above conditional probability is computed as follows: $P\{X > a + b | X > a\} = \frac{P\{X > a + b\}}{P\{X > a\}} = \left(\frac{x_m}{a + b}\right)^{\alpha} \cdot \left(\frac{a}{x_m}\right)^{\alpha} = \left(\frac{a}{a + b}\right)^{\alpha}$ The distribution is not memoryless, since the above is not equal to $P\{X > b\} = \left(\frac{x_m}{b}\right)^{\alpha}$. When $a \to +\infty$, the conditional probability approaches 1. This means that if you have exceeded a very large value, then the probability that you will exceed an even larger one approaches 1.

Exercise 2 - Solution

1) The CTMC is the following. Note that the state of the system cannot be described using only the number of jobs, since the fact that the system is *decrypting* or *serving* jobs is also relevant. Starred states are those when the system is decrypting. Furthermore, you can observe that when $\alpha \to \infty$ (i.e., decryption is much faster than the rate of arrival or service) the two states in the same column collapse to one, and the system becomes an M/M/1.



2) Using local equilibrium equations, one can easily check that $p_j \cdot \lambda = p_{j+1} \cdot \mu$, hence $p_j = \left(\frac{\lambda}{\mu}\right)^j \cdot p_0$, $j \ge 0$. Moreover, using global equilibrium equation around the decrypting states, one obtains $p_j^* \cdot \alpha = p_{j-1} \cdot \lambda$, $j \ge 1$, hence $p_j^* = \frac{\mu}{\alpha} \cdot \left(\frac{\lambda}{\mu}\right)^j \cdot p_0$, $j \ge 1$. The normalization condition then reads $p_0 \left[\sum_{j=0}^{+\infty} \left(\frac{\lambda}{\mu}\right)^j + \frac{\mu}{\alpha} \cdot \sum_{j=1}^{+\infty} \left(\frac{\lambda}{\mu}\right)^j \right] = 1$. Call $\rho = \frac{\lambda}{\mu}$, the stability condition is then $\rho < 1$, and it is independent of α . This is as expected, since during started states the system is neither accepting nor serving jobs, so the decryption rate cannot have any influence on stability. After a few straightforward computations, one obtains:

$$p_j = (1 - \rho) \cdot \rho^j \cdot \frac{\alpha}{\lambda + \alpha}, \ j \ge 0.$$
$$p_j^* = (1 - \rho) \cdot \rho^j \cdot \frac{\mu}{\lambda + \alpha}, \ j \ge 1.$$

3) The average arrival rate in the system is $\overline{\lambda} = \lambda \cdot \sum_{j=0}^{+\infty} p_j + 0 \cdot \sum_{j=0}^{+\infty} p_j^* = \lambda \cdot \frac{\alpha}{\lambda + \alpha}$. The tagged arrival probabilities are $r_j = p_j \cdot \frac{\lambda}{\overline{\lambda}} = (1 - \rho) \cdot \rho^j$, $r_j^* = 0$.

4) The mean number of jobs in the system (which includes a job being decrypted, if there is one) is:

$$E[N] = \sum_{j=1}^{+\infty} j \cdot \left(p_j + p_j^*\right) = \frac{\mu + \alpha}{\lambda + \alpha} \cdot \sum_{j=1}^{+\infty} j \cdot (1 - \rho) \cdot \rho^j = \frac{\mu + \alpha}{\lambda + \alpha} \cdot \frac{\rho}{1 - \rho}$$

The mean traversal time for a job can be found by applying Little's Law, and it is:

$$E[R] = \frac{E[N]}{\overline{\lambda}} = \frac{\mu + \alpha}{\lambda + \alpha} \cdot \frac{\rho}{1 - \rho} \cdot \frac{\lambda + \alpha}{\lambda \cdot \alpha} = \frac{\mu + \alpha}{\alpha} \cdot \frac{1}{\mu - \lambda}$$