




# Perfect Forward Security

Security in Networked Computing Systems

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## PKE-based Key Ex

|   |                                    |  |
|---|------------------------------------|--|
| <p><b>A</b><br/><b>(pubK<sub>B</sub>)</b></p> |                                    | <p><b>B</b><br/><b>(privK<sub>B</sub>, pubK<sub>B</sub>)</b></p> |
| <p>K ← random()</p>                           | <p>CT = E(pubK<sub>B</sub>, K)</p> | <p>K = D(privK<sub>B</sub>, CT)</p>                              |
|   | <p>E(K, session)</p>               |  |
| <p>Delete K</p>                               |                                    | <p>Delete K</p>  |

- Private key **privK<sub>B</sub>** is a *long-term* secret
- Key **K** is the *session* key
- SSL/TLS employs a similar scheme

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## The problem




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- The adversary records CTs of the session
- If the adversary compromises  $\text{priv}K_B$  then (s)he can recover K from CT
- Then, the adversary decrypts the session and violates secrecy
- The long-term secret becomes a single-point of failure

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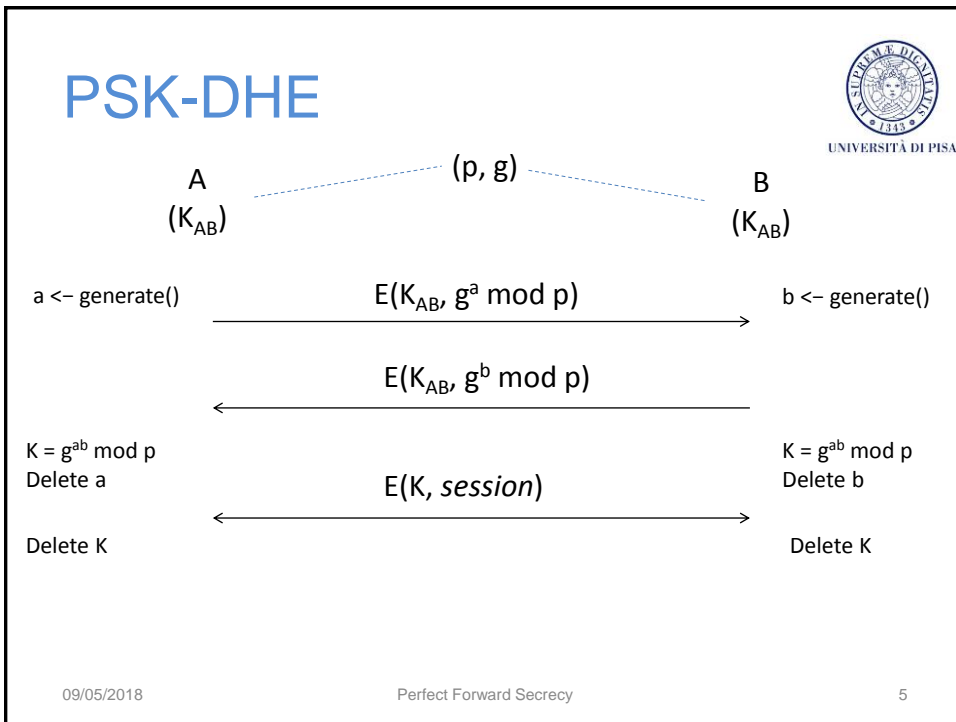
## Perfect Forward Secrecy




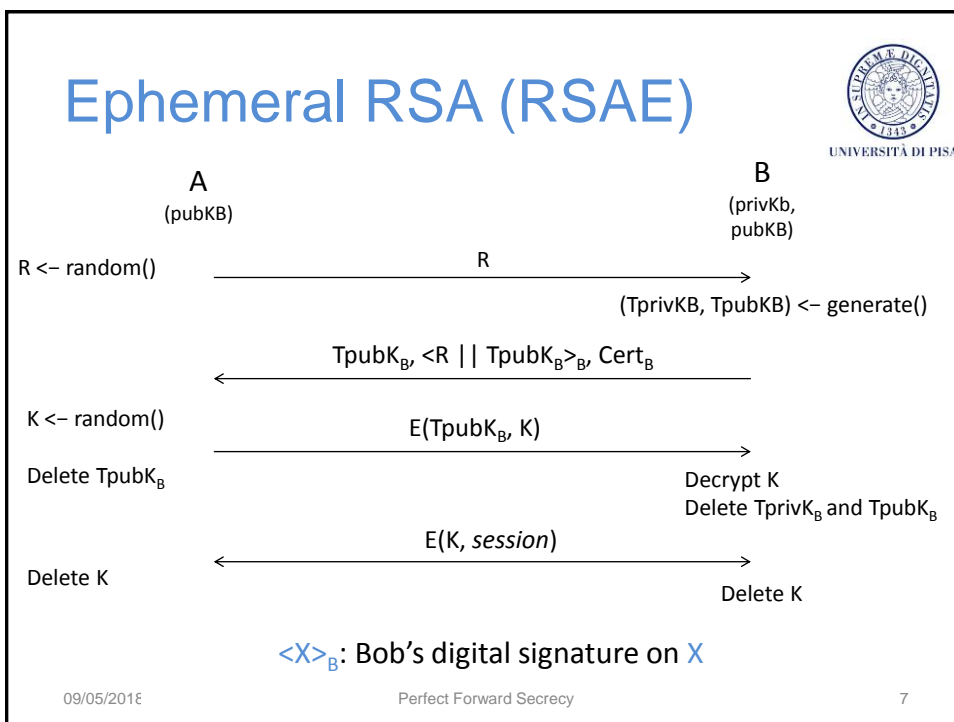
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- **(DEF) PFS:** Disclosure of long-term secret keying material does not compromise the secrecy of the exchanged keys from earlier runs
- PKE, and in particular DH, makes it possible to achieve this requirement


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- # PSK-DHE
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- Pre-Shared Key Ephemeral Diffie-Hellman
  - Ephemeral Diffie-Hellman
    - Keys  $a$  and  $b$  are ephemeral (one-time per-session or per message)
  - Once  $a$  and  $b$  (and  $K$ ) have been deleted there is no way to recover  $K$ , and thus the session, even if the long-term private  $K_{ab}$  is compromised: neither A nor B can
    - Even though private  $K_{ab}$  is compromised, you still have to solve the DLP
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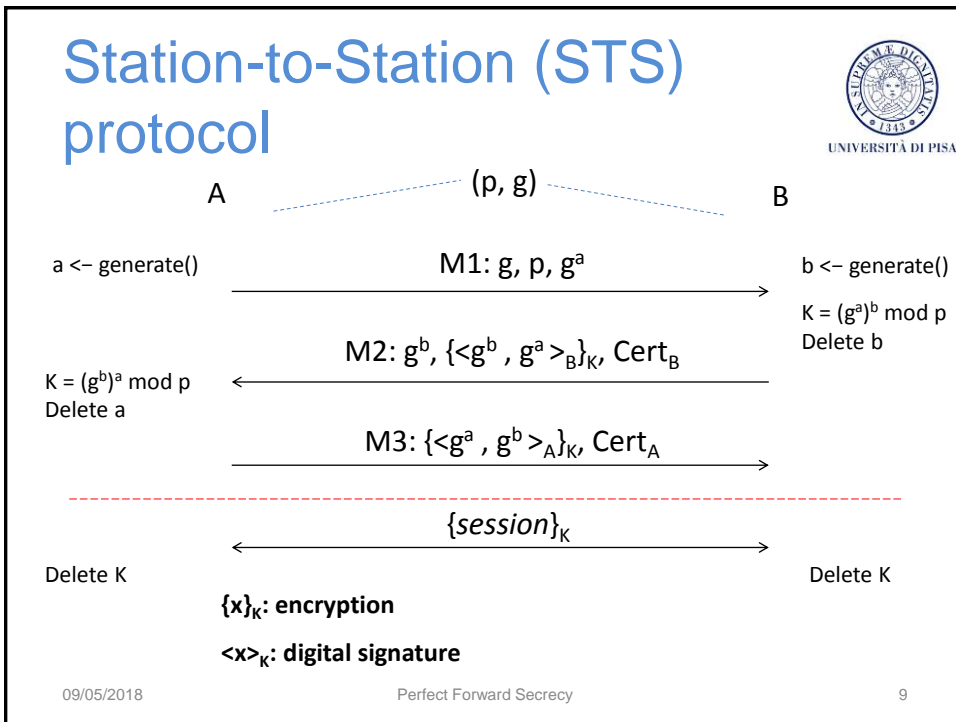
## Direct Authentication




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- **(DEF) Direct Authentication** means to prove the peer the knowledge of K
  - If a KeyEx protocol does not fulfil direct authentication, this authentication is achieved at the first application message
  - DA is said *key confirmation* in the BAN parlance
- DHE and RSAE don't fulfil direct authentication
  - Until E(K, session)
- **Station-To-Station (STS) Protocol** fulfils direct authentication

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- ## Misc
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- **CONS**
    - PFS requires more computation
    - Crypto-(co)processors do not support PFS (for the moment)
  - **Who uses PFS**
    - Whatsapp, Twitter, IOS9, Google
    - (EC)DHE is part of SSL/TLS cipher suite
  - **SSL Quality Test**
    - <https://www.ssllabs.com/ssltest>
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