EXERCISE 1

Construct a B+-tree for the following set of values

Assume that the tree is initially empty and values are inserted in ascending order.

1) Construct B+-trees for the cases where the number *m* of pointers that will fit a node is as follows: a. Four

b. Seven

2) Shows the form of the B+-tree after each operation of the sequence: Insert 9; Insert 10; Insert 8; Insert 6; Insert 1; Insert 4 for the case *m*=4.

Solution

Point 1)



B+-tree:

Not a root or a leaf node: $\lceil m/2 \rceil \le p \le m$, where p is a pointer. Leaf node: $\lceil (m-1)/2 \rceil \le k \le m-1$, where k is a key. Root: $p \ge 2$ if it is not a leaf

 $0 \le k \le (m-1)$ if it is a leaf.

a) <i>m</i> =4	Internal node: 2<=p<=4	Root: $p \ge 2$ if it is not a leaf;
	Leaf node: 2<=k<=3	$0 \le k \le 3$ if it is a leaf

B+tree: Insert 2, 3, 5

B+tree: Insert 7 (leaf split)



Split a leaf node:

Let K_1, \ldots, K_m be the set of keys in the ascending order (i.e. 2, 3, 5, 7)

- Node N: $K_1, \ldots, K_{\lceil m/2 \rceil 1}, K_{\lceil m/2 \rceil}$ (i.e. 2, 3)
- Node N': $K_{\lceil m/2 \rceil+1}, ..., K_m$ (i.e. 5, 7)
- Insert $K_{\lceil m/2 \rceil+1}$ into the parent node (i.e. 5)

B+tree: Insert 11, 17, 19, 23, 29 11 | 19 5 5 | 7 | 11 | 17 | 19 | 23 | 29 2 | 3 |

B+tree: Insert 31 (leaf node split + non leaf node split)



Split an internal node N:

Let K_1, \ldots, K_m be the set of keys in the ascending order (i.e. 5, 11, 19, 29)

- Node N: $K_1, \ldots, K_{\lceil m/2 \rceil 1}$ (i.e. 5)
- Node N': $K_{\lceil m/2 \rceil+1}$, ..., K_m (i.e. 19, 29)
- Insert $K \upharpoonright_{m/2}$ into the parent node of N (i.e. 11)

Note that since values are inserted in ascending order, leaves have the minimum number of keys except the last leaf.

b) Seven.

m=7

Intermediate node : 4<=p<=7 Leaf node: $3 \le k \le 6$ Root: $p \ge 2$ if it is not a leaf; $0 \le k \le 6$ if it is a leaf

B+tree: Insert 2, 3, 5, 7, 11, 17



B+tree: Insert 19 (leaf split)







B+tree: Insert 10 (leaf split)





B+tree: Insert 6 (leaf split)





B+tree: Insert 4 (leaf node split + non leaf node split)



EXERCISE 2 (leaf merge, non leaf merge, leaf keys redistribution)

For the following B+-tree (m = 5) show the form of the tree after each of the of operations of the sequence:

Delete 17; Delete 20; Delete 34

What is the cost in terms of block transfers for each operation?



Solution

m=5 Intermediate node : $3 \le p \le 5$ Root: $p \ge 2$ if it is not a leaf ; $0 \le k \le 4$ if it is a leaf Leaf node: 2<=k<=4

Delete k:

Find the leaf node that contains k; Delete k from the node; If the node has too few entries

- 1) merge nodes (if possible)
- 2) otherwise redistribute keys

B+tree: Delete 17



Cost = 3 read + 1 write = 4



Leaf merge:

Let N' be the predecessor; let N be the successor Let k' be the value between the two nodes N' and N in the parent Append all keys to N' Delete (k, pointer to N) from the parent Delete N



Non leaf merge:

Node Q too few pointers Let Q' predecessor and Q successor Let K' be the value between the two nodes in parent of Q Append K' and all pointers and values of Q to Q' Delete (K', pointer) from the parent



Root has only one child. Root can be deleted.



Cost: 5 read + 2 write = 7

Delete 34: keys redistribution (leaf)



The node N has too few values

Let N' be the previous or next child of parent of N

Let K' the value between N and N' in the parent

Entries of N and N' cannot fit in a single node.

We apply redistribution of keys

N borrows an entry from N' (assume N' predecessor of N)

Let j such that (Pj, Kj) is the last (pointer, value) in N'

Remove (Pj, Kj) from N'

Insert (Pj, Kj) as first value in N

Replace K' by Kj in the parent



Cost = 3 read + 3 write = 6

EXERCISE 3 (non leaf keys redistribuiton)

Show the form of the B+-tree (m=5) after the operation Delete 20 What is the cost of the operation?



Solution

Intermediate node : 3<=p<=5 Leaf node: 2<=k<=4 m=5 Root: $p \ge 2$ if it is not a leaf; $0 \le k \le 4$ if it is a leaf

Delete 20



The node N has too few values Merge between N and N'



The node Q has too few pointers

No merge with previous or net child of parent of Q Redistribution of keys Q borrows an entry from Q' Let j be such that (Kj-1, pj) is the last value pointer in Q' Let K' the value between Q and Q' in parent of Q Insert (pj, K') as first value of Q Remove (Kj-1, pj) from Q' Replace K' with Kj-1 in parent of Q



Cost = 5 read + 4 write = 9